

Specifying and Checking Java using CSP

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Content

- Why CSP?
- CSP-OZ to Java
- jassda
- Differences to CSP_M
- Conclusion and Future Work

Why CSP?

- Project ForMooS: **F**ormal **M**ethods in **o**bject oriented **S**oftware Engineering
- base: CSP-OZ [C. Fischer 2000]
 - *Communicating Sequential Processes* [Hoare '85] +
 - *Object-Z* [Smith 2000]
- goal: UML (subset) → CSP-OZ → Java

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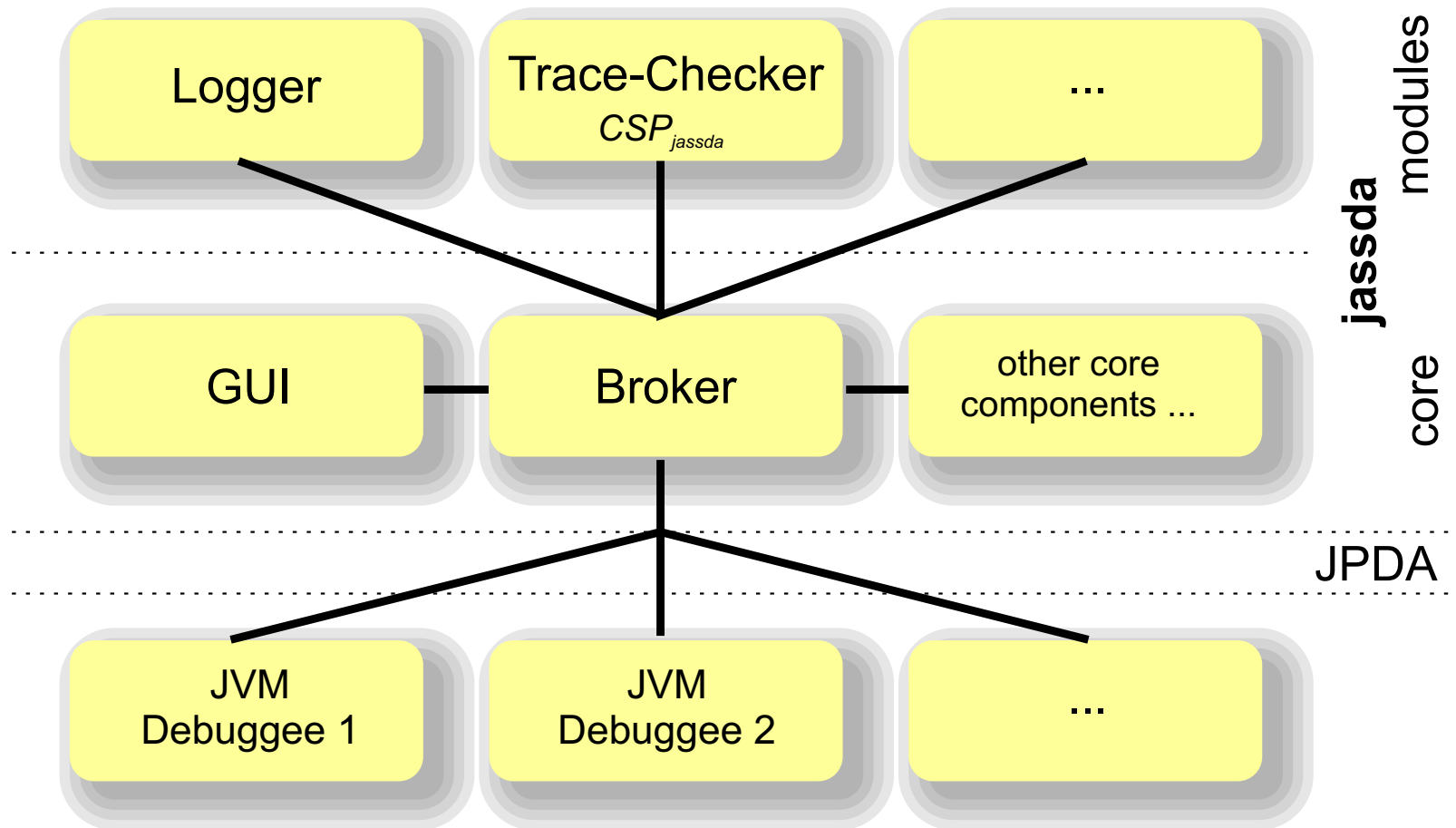
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 2. Trace Assertions
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skip CSP-OZ

jassda Overview

- ➔ Jass: **J**ava with **ass**ertions
- ➔ jassda: **Jass** **D**ebug **A**rchitecture
- ➔ Debug on Byte-code level
- ➔ modular structure (framework for debugging)

jassda Architecture



Serialised Event Stream

- Using JDI means getting a **serialised** stream of events.
- serialised stream = trace
- CSP trace semantics
 - check: (trace of current run) \in CSP trace semantics
 - “checking operational”:
do not expand the trace semantics
describe further trace(s) by processes (for every step)

Possible *events*: Everything that can stop the VM via JDI.
CSP_{jassda} basic events: (analogous to Design by Contract clauses)

- Method entry point `begin`
- Normal method termination `end`
- Exceptional method termination `exception`

Event Properties



- ➡ Virtual Machine (Debuggee)
- ➡ JDI

Event Properties



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- ➡ class
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Filtering of events through properties by Java classes
(handler class)

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- ➡ definition of event sets

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eventset myset = { handler="..."  
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- operations on sets: intersection, union

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▣ Variable binding in CSP_{jassda} :

```
main() { in?x:[arg0] -> out!x -> main() }
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Internal Choice (CSP_M)

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- so: use trace semantics

$$traces(a \rightarrow P \sqcap b \rightarrow Q) = traces(a \rightarrow P \square b \rightarrow Q)$$

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Hello World



```
eventset helloWorld
  { handler="jass.debugger.jdi.eventset.GenericSet",
    class  ="HelloWorld" }
eventset start { eventtype="begin" method="start" }
eventset stop  { eventtype="begin" method="stop" }

main() {
  ||x:[instance] @ helloWorldProc(x)
}
helloWorldProc(x) {
  helloWorld.start.x ->
    helloWorld.stop.x -> helloWorldProc(x)
}
```

Conclusion



➡ *CSP_{jassda}*:

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- ⇒ extendable through modular structure

Future Work

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Many thanks to Mark Brörkens